

-NoValue-

syn.1 Semantic Notions

mvl:syn:sem:
sec Suppose a many-valued logic $\mathbf{L}$ is given by a matrix. Then we can define the usual semantic notions for $\mathbf{L}$.

- Definição syn.1.** 1. A fórmula φ is *satisfiable* if for some $\mathbf{v}$, $\mathbf{v} \models \varphi$; it is *unsatisfiable* if for no $\mathbf{v}$, $\mathbf{v} \models \varphi$;
2. A fórmula φ is a *tautology* if $\mathbf{v} \models \varphi$ for all *valorações* $\mathbf{v}$;
3. If Γ is a set of *fórmulas*, $\Gamma \models \varphi$ (“ Γ entails φ ”) if and only if $\mathbf{v} \models \varphi$ for every *valoração* $\mathbf{v}$ for which $\mathbf{v} \models \Gamma$.
4. If Γ is a set of *fórmulas*, Γ is *satisfiable* if there is a *valoração* $\mathbf{v}$ for which $\mathbf{v} \models \Gamma$, and Γ is *unsatisfiable* otherwise.

We have some of the same facts for these notions as we do for the case of classical logic:

mvl:syn:sem:
prop:semanticalfacts **Proposição syn.2.**

1. φ is a tautology if and only if $\emptyset \models \varphi$;
2. If Γ is satisfiable then every finite subset of Γ is also satisfiable;
3. Monotonicity: if $\Gamma \subseteq \Delta$ and $\Gamma \models \varphi$ then also $\Delta \models \varphi$;
4. Transitivity: if $\Gamma \models \varphi$ and $\Delta \cup \{\varphi\} \models \psi$ then $\Gamma \cup \Delta \models \psi$;

Prova. Exercise. □

Problema syn.1. Prove *Proposição syn.2*

In classical logic we can connect entailment and the conditional. For instance, we have the validity of *modus ponens*: If $\Gamma \models \varphi$ and $\Gamma \models \varphi \rightarrow \psi$ then $\Gamma \models \psi$. Another important relationship between $\models$ and $\rightarrow$ in classical logic is the semantic deduction theorem: $\Gamma \models \varphi \rightarrow \psi$ if and only if $\Gamma \cup \{\varphi\} \models \psi$. These results *do not* always hold in many-valued logics. Whether they do depends on the truth function $\widetilde{\rightarrow}$.

Créditos das Imagens

Referências Bibliográficas