

-NoValue-

thy.1 The s - m - n Theorem

cmp:thy:smn:sec The next theorem is known as the “ s - m - n theorem,” for a reason that will be explanation clear in a moment. The hard part is understanding just what the theorem says; once you understand the statement, it will seem fairly obvious.

cmp:thy:smn:thm:s-m-n **Teorema thy.1.** *For each pair of natural numbers n and m , there is a primitive recursive function s_n^m such that for every sequence $e, a_0, \dots, a_{m-1}, y_0, \dots, y_{n-1}$, we have*

$$\varphi_{s_n^m(e, a_0, \dots, a_{m-1})}^n(y_0, \dots, y_{n-1}) \simeq \varphi_e^{m+n}(a_0, \dots, a_{m-1}, y_0, \dots, y_{n-1}).$$

It is helpful to think of s_n^m as acting on *programs*. That is, s_n^m takes a explanation program e for an $(m+n)$ -ary function, as well as fixed inputs $a_0, \dots, a_{m-1}$; and it returns a program $s_n^m(x, a_0, \dots, a_{m-1})$ for the n -ary function of the remaining arguments. If you think of x as the description of a Turing machine, then $s_n^m(e, a_0, \dots, a_{m-1})$ is the Turing machine that, on input $y_0, \dots, y_{n-1}$, prepends $a_0, \dots, a_{m-1}$ to the input string, and runs e . Each s_n^m is then just a primitive recursive function that finds a code for the appropriate Turing machine.

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Referências Bibliográficas