

## tab.1 Countermodels from Tableaux

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sec

The proof of the completeness theorem doesn't just show that if  $\models \varphi$  then  $\vdash \varphi$ , [explanation](#) it also gives us a method for constructing countermodels to  $\varphi$  if  $\not\models A$ . In the case of  $\mathbf{K}$ , this method constitutes a *decision procedure*. For suppose  $\not\models \varphi$ . Then the proof of ?? gives a method for constructing a complete [tableau](#). The method in fact always terminates. The propositional rules for  $\mathbf{K}$  only add prefixed [formulas](#) of lower complexity, i.e., each propositional rule need only be applied once on a branch for any signed formula  $\sigma S \varphi$ . New prefixes are only generated by the  $\Box\mathbb{F}$  and  $\Diamond\mathbb{T}$  rules, and also only have to be applied once (and produce a single new prefix).  $\Box\mathbb{T}$  and  $\Diamond\mathbb{F}$  have to be applied potentially multiple times, but only once per prefix, and only finitely many new prefixes are generated. So the construction either results in a closed branch or a complete branch after finitely many stages.

Once a tableau with an open complete branch is constructed, the proof of ?? gives us an explicit model that satisfies the original set of prefixed [formulas](#). So not only is it the case that if  $\Gamma \models \varphi$ , then a closed [tableau](#) exists and  $\Gamma \vdash \varphi$ , if we look for the closed [tableau](#) in the right way and end up with a “complete” [tableau](#), we'll not only know that  $\Gamma \not\models \varphi$  but actually be able to construct a countermodel.

**Example tab.1.** We know that  $\not\models \Box(p \vee q) \rightarrow (\Box p \vee \Box q)$ . The construction of a tableau begins with:

1.	$1\mathbb{F} \Box(p \vee q) \rightarrow (\Box p \vee \Box q) \checkmark$	Assumption
2.	$1\mathbb{T} \Box(p \vee q)$	$\rightarrow\mathbb{F} 1$
3.	$1\mathbb{F} \Box p \vee \Box q \checkmark$	$\rightarrow\mathbb{F} 1$
4.	$1\mathbb{F} \Box p \checkmark$	$\vee\mathbb{F} 3$
5.	$1\mathbb{F} \Box q \checkmark$	$\vee\mathbb{F} 3$
6.	$1.1\mathbb{F} p \checkmark$	$\Box\mathbb{F} 4$
7.	$1.2\mathbb{F} q \checkmark$	$\Box\mathbb{F} 5$

The [tableau](#) is of course not finished yet. In the next step, we consider the only line without a checkmark: the prefixed [formula](#)  $1\mathbb{T}\Box(p \vee q)$  on line 2. The construction of the closed tableau says to apply the  $\Box\mathbb{T}$  rule for every prefix used on the branch, i.e., for both 1.1 and 1.2:

1.	$1\mathbb{F} \Box(p \vee q) \rightarrow (\Box p \vee \Box q) \checkmark$	Assumption
2.	$1\mathbb{T} \Box(p \vee q)$	$\rightarrow\mathbb{F} 1$
3.	$1\mathbb{F} \Box p \vee \Box q \checkmark$	$\rightarrow\mathbb{F} 1$
4.	$1\mathbb{F} \Box p \checkmark$	$\vee\mathbb{F} 3$
5.	$1\mathbb{F} \Box q \checkmark$	$\vee\mathbb{F} 3$
6.	$1.1\mathbb{F} p \checkmark$	$\Box\mathbb{F} 4$
7.	$1.2\mathbb{F} q \checkmark$	$\Box\mathbb{F} 5$
8.	$1.1\mathbb{T} p \vee q$	$\Box\mathbb{T} 2$
9.	$1.2\mathbb{T} p \vee q$	$\Box\mathbb{T} 2$

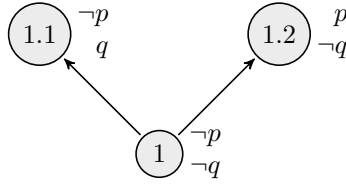


Figure 1: A countermodel to  $\Box(p \vee q) \rightarrow (\Box p \vee \Box q)$ .

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fig:counter-Box

Now lines 2, 8, and 9, don't have checkmarks. But no new prefix has been added, so we apply  $\forall\mathbb{T}$  to lines 8 and 9, on all resulting branches (as long as they don't close):

1.	1F $\Box(p \vee q) \rightarrow (\Box p \vee \Box q) \checkmark$	Assumption
2.	1T $\Box(p \vee q)$	$\rightarrow\mathbb{F}$ 1
3.	1F $\Box p \vee \Box q \checkmark$	$\rightarrow\mathbb{F}$ 1
4.	1F $\Box p \checkmark$	$\forall\mathbb{F}$ 3
5.	1F $\Box q \checkmark$	$\forall\mathbb{F}$ 3
6.	1.1F $p \checkmark$	$\Box\mathbb{F}$ 4
7.	1.2F $q \checkmark$	$\Box\mathbb{F}$ 5
8.	1.1T $p \vee q \checkmark$	$\Box\mathbb{T}$ 2
9.	1.2T $p \vee q \checkmark$	$\Box\mathbb{T}$ 2
<div style="display: flex; justify-content: space-around; width: 100%;"> <div style="text-align: center;"> <math>\swarrow</math>            10. 1.1T <math>p \checkmark</math>  <math>\otimes</math> </div> <div style="text-align: center;"> <math>\searrow</math>            1.1T <math>q \checkmark</math>  <math>\swarrow</math>            11. 1.2T <math>p \checkmark</math>    1.2T <math>q \checkmark</math>  <math>\otimes</math> </div> </div>		
		$\forall\mathbb{T}$ 8
		$\forall\mathbb{T}$ 9

There is one remaining open branch, and it is complete. From it we define the model with worlds  $W = \{1, 1.1, 1.2\}$  (the only prefixes appearing on the open branch), the accessibility relation  $R = \{\langle 1, 1.1 \rangle, \langle 1, 1.2 \rangle\}$ , and the assignment  $V(p) = \{1.2\}$  (because line 11 contains  $1.2\mathbb{T}p$ ) and  $V(q) = \{1.1\}$  (because line 10 contains  $1.1\mathbb{T}q$ ). The model is pictured in [Figure 1](#), and you can verify that it is a countermodel to  $\Box(p \vee q) \rightarrow (\Box p \vee \Box q)$ .

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## Bibliography