

## thr.1 Gödel logics

mvl:thr:god:sec Kurt Gödel introduced a sequence of  $n$ -valued logics that each contain all formulas valid in intuitionistic logic, and are contained in classical logic. Here is the first interesting one:

mvl:thr:god:defn:goedel **Definition thr.1.** *3-valued Gödel logic*  $\mathbf{G}$  is defined using the matrix:

1. The standard propositional language  $\mathcal{L}_0$  with  $\perp, \neg, \wedge, \vee, \rightarrow$ .
2. The set of truth values  $V = \{\mathbb{T}, \mathbb{U}, \mathbb{F}\}$ .
3.  $\mathbb{T}$  is the only designated value, i.e.,  $V^+ = \{\mathbb{T}\}$ .
4. For  $\perp$ , we have  $\tilde{\perp} = \mathbb{F}$ . Truth functions for the remaining connectives are given by the following tables:

$\tilde{\neg}_{\mathbf{G}}$		$\tilde{\wedge}_{\mathbf{G}}$	$\mathbb{T}$	$\mathbb{U}$	$\mathbb{F}$
$\mathbb{T}$	$\mathbb{F}$	$\mathbb{T}$	$\mathbb{T}$	$\mathbb{U}$	$\mathbb{F}$
$\mathbb{U}$	$\mathbb{F}$	$\mathbb{U}$	$\mathbb{U}$	$\mathbb{U}$	$\mathbb{F}$
$\mathbb{F}$	$\mathbb{T}$	$\mathbb{F}$	$\mathbb{F}$	$\mathbb{F}$	$\mathbb{F}$

  

$\tilde{\vee}_{\mathbf{G}}$	$\mathbb{T}$	$\mathbb{U}$	$\mathbb{F}$	$\tilde{\rightarrow}_{\mathbf{G}}$	$\mathbb{T}$	$\mathbb{U}$	$\mathbb{F}$
$\mathbb{T}$	$\mathbb{T}$	$\mathbb{T}$	$\mathbb{T}$	$\mathbb{T}$	$\mathbb{T}$	$\mathbb{U}$	$\mathbb{F}$
$\mathbb{U}$	$\mathbb{T}$	$\mathbb{U}$	$\mathbb{U}$	$\mathbb{U}$	$\mathbb{T}$	$\mathbb{T}$	$\mathbb{F}$
$\mathbb{F}$	$\mathbb{T}$	$\mathbb{U}$	$\mathbb{F}$	$\mathbb{F}$	$\mathbb{T}$	$\mathbb{T}$	$\mathbb{T}$

You'll notice that the truth tables for  $\wedge$  and  $\vee$  are the same as in Łukasiewicz and strong Kleene logic, but the truth tables for  $\neg$  and  $\rightarrow$  differ for each. In Gödel logic,  $\tilde{\neg}(\mathbb{U}) = \mathbb{F}$ . In contrast to Łukasiewicz logic and Kleene logic,  $\tilde{\rightarrow}(\mathbb{U}, \mathbb{F}) = \mathbb{F}$ ; in contrast to Kleene logic (but as in Łukasiewicz logic),  $\tilde{\rightarrow}(\mathbb{U}, \mathbb{U}) = \mathbb{T}$ .

As the connection to intuitionistic logic alluded to above suggests,  $\mathbf{G}_3$  is close to intuitionistic logic. All intuitionistic truths are tautologies in  $\mathbf{G}_3$ , and many classical tautologies that are not valid intuitionistically also fail to be tautologies in  $\mathbf{G}_3$ . For instance, the following are not tautologies:

$$\begin{array}{ll}
 p \vee \neg p & (p \rightarrow q) \rightarrow (\neg p \vee q) \\
 \neg \neg p \rightarrow p & \neg(p \wedge q) \rightarrow (\neg p \vee \neg q) \\
 & ((p \rightarrow q) \rightarrow p) \rightarrow p
 \end{array}$$

However, not every tautology of  $\mathbf{G}_3$  is also intuitionistically valid, e.g.,  $(p \rightarrow q) \vee (q \rightarrow p)$ .

**Problem thr.1.** Give a truth table to show that  $(p \rightarrow q) \vee (q \rightarrow p)$  is a tautology of  $\mathbf{G}_3$ .

**Problem thr.2.** Give truth tables that show that the following are not tautologies of  $\mathbf{G}_3$

$$(p \rightarrow q) \rightarrow (\neg p \vee q)$$
$$\neg(p \wedge q) \rightarrow (\neg p \vee \neg q)$$
$$((p \rightarrow q) \rightarrow p) \rightarrow p$$

**Problem thr.3.** Which of the following relations hold in Gödel logic? Give a truth table for each.

1.  $p, p \rightarrow q \models q$
2.  $p \vee q, \neg p \models q$
3.  $p \wedge q \models p$
4.  $p \models p \wedge p$
5.  $p \models p \vee q$

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**Bibliography**