

## syn.1 Semantic Notions

mvl:syn:sem:  
sec Suppose a many-valued logic  $\mathbf{L}$  is given by a matrix. Then we can define the usual semantic notions for  $\mathbf{L}$ .

- Definition syn.1.**
1. A formula  $\varphi$  is *satisfiable* if for some  $\mathbf{v}$ ,  $\mathbf{v} \models \varphi$ ; it is *unsatisfiable* if for no  $\mathbf{v}$ ,  $\mathbf{v} \models \varphi$ ;
  2. A formula  $\varphi$  is a *tautology* if  $\mathbf{v} \models \varphi$  for all valuations  $v$ ;
  3. If  $\Gamma$  is a set of formulas,  $\Gamma \models \varphi$  (“ $\Gamma$  entails  $\varphi$ ”) if and only if  $\mathbf{v} \models \varphi$  for every valuation  $\mathbf{v}$  for which  $\mathbf{v} \models \Gamma$ .
  4. If  $\Gamma$  is a set of formulas,  $\Gamma$  is *satisfiable* if there is a valuation  $\mathbf{v}$  for which  $\mathbf{v} \models \Gamma$ , and  $\Gamma$  is *unsatisfiable* otherwise.

We have some of the same facts for these notions as we do for the case of classical logic:

mvl:syn:sem:  
prop:semanticalfacts **Proposition syn.2.**

1.  $\varphi$  is a tautology if and only if  $\emptyset \models \varphi$ ;
2. If  $\Gamma$  is satisfiable then every finite subset of  $\Gamma$  is also satisfiable;
3. *Monotony:* if  $\Gamma \subseteq \Delta$  and  $\Gamma \models \varphi$  then also  $\Delta \models \varphi$ ;
4. *Transitivity:* if  $\Gamma \models \varphi$  and  $\Delta \cup \{\varphi\} \models \psi$  then  $\Gamma \cup \Delta \models \psi$ ;

mvl:syn:sem:  
def:Monotony

mvl:syn:sem:  
def:Cut

*Proof.* Exercise. □

**Problem syn.1.** Prove [Proposition syn.2](#)

In classical logic we can connect entailment and the conditional. For instance, we have the validity of *modus ponens*: If  $\Gamma \models \varphi$  and  $\Gamma \models \varphi \rightarrow \psi$  then  $\Gamma \models \psi$ . Another important relationship between  $\models$  and  $\rightarrow$  in classical logic is the semantic deduction theorem:  $\Gamma \models \varphi \rightarrow \psi$  if and only if  $\Gamma \cup \{\varphi\} \models \psi$ . These results *do not* always hold in many-valued logics. Whether they do depends on the truth function  $\widetilde{\rightarrow}$ .

## Photo Credits

## Bibliography