int.1   The Syntax of the Lambda Calculus

One starts with a sequence of variables x, y, z, ... and some constant symbols a, b, c, ... The set of terms is defined inductively, as follows:

1. Each variable is a term.

2. Each constant is a term.

3. If M and N are terms, so is (MN).

4. If M is a term and x is a variable, then $(\lambda x.\, M)$ is a term.

The system without any constants at all is called the *pure* lambda calculus. We'll mainly be working in the pure λ-calculus, so all lowerce letters will stand for variables. We use uppercase letters (M, N, etc.) to stand for terms of the λ-calculus.

We will follow a few notational conventions:

Convention 1.   1. When parentheses are left out, application takes place from left to right. For example, if M, N, P, and Q are terms, then $MNPQ$ abbreviates $(((MN)P)Q)$.

2. Again, when parentheses are left out, lambda abstraction is to be given the widest scope possible. From example, $\lambda x.\, MNP$ is read $(\lambda x.\, ((MN)P))$.

3. A lambda can be used to abstract multiple variables. For example, $\lambda xyz.\, M$ is short for $\lambda x.\, \lambda y.\, \lambda z.\, M$.

For example,

$$\lambda xy.\, xxyx\lambda z.\, xz$$

abbreviates

$$\lambda x.\, \lambda y.\, ((((xx)y)x)(\lambda z.\, (xz))).$$

You should memorize these conventions. They will drive you crazy at first, but you will get used to them, and after a while they will drive you less crazy than having to deal with a morass of parentheses.

Two terms that differ only in the names of the bound variables are called α-equivalent; for example, $\lambda x.\, x$ and $\lambda y.\, y$. It will be convenient to think of these as being the "same" term; in other words, when we say that M and N are the same, we also mean "up to renamings of the bound variables." Variables that are in the scope of a λ are called "bound", while others are called "free." There are no free variables in the previous example; but in

$$(\lambda z.\, yz)x$$

y and x are free, and z is bound.

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Bibliography