

## int.1 The Syntax of the Lambda Calculus

lam:int:syn:  
sec One starts with a sequence of variables  $x, y, z, \dots$  and some constant symbols  $a, b, c, \dots$ . The set of terms is defined inductively, as follows:

1. Each variable is a term.
2. Each constant is a term.
3. If  $M$  and  $N$  are terms, so is  $(MN)$ .
4. If  $M$  is a term and  $x$  is a variable, then  $(\lambda x. M)$  is a term.

The system without any constants at all is called the *pure* lambda calculus. We'll mainly be working in the pure  $\lambda$ -calculus, so all lowercase letters will stand for variables. We use uppercase letters ( $M, N$ , etc.) to stand for terms of the  $\lambda$ -calculus.

We will follow a few notational conventions:

- Convention 1.*
1. When parentheses are left out, application takes place from left to right. For example, if  $M, N, P$ , and  $Q$  are terms, then  $MNPQ$  abbreviates  $((MN)P)Q$ .
  2. Again, when parentheses are left out, lambda abstraction is to be given the widest scope possible. For example,  $\lambda x. MNP$  is read  $(\lambda x. ((MN)P))$ .
  3. A lambda can be used to abstract multiple variables. For example,  $\lambda xyz. M$  is short for  $\lambda x. \lambda y. \lambda z. M$ .

For example,

$$\lambda xy. xxyx \lambda z. xz$$

abbreviates

$$\lambda x. \lambda y. (((xxy)x)(\lambda z. (xz))).$$

You should memorize these conventions. They will drive you crazy at first, but you will get used to them, and after a while they will drive you less crazy than having to deal with a morass of parentheses.

Two terms that differ only in the names of the bound variables are called  $\alpha$ -equivalent; for example,  $\lambda x. x$  and  $\lambda y. y$ . It will be convenient to think of these as being the “same” term; in other words, when we say that  $M$  and  $N$  are the same, we also mean “up to renamings of the bound variables.” Variables that are in the scope of a  $\lambda$  are called “bound”, while others are called “free.” There are no free variables in the previous example; but in

$$(\lambda z. yz)x$$

$y$  and  $x$  are free, and  $z$  is bound.

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**Bibliography**