

## sc.1 The Canonical Model

int:sc:mod:  
sec The worlds in our model will be finite sequences  $\sigma$  of natural numbers, i.e.,  $\sigma \in \mathbb{N}^*$ . Note that  $\mathbb{N}^*$  is inductively defined by:

1.  $\Lambda \in \mathbb{N}^*$ .
2. If  $\sigma \in \mathbb{N}^*$  and  $n \in \mathbb{N}$ , then  $\sigma.n \in \mathbb{N}^*$  (where  $\sigma.n$  is  $\sigma \frown \langle n \rangle$  and  $\sigma \frown \sigma'$  is the concatenation of  $\sigma$  and  $\sigma'$ ).
3. Nothing else is in  $\mathbb{N}^*$ .

So we can use  $\mathbb{N}^*$  to give inductive definitions.

Let  $\langle \psi_1, \chi_1 \rangle, \langle \psi_2, \chi_2 \rangle, \dots$ , be an enumeration of all pairs of **formulas**. Given a set of **formulas**  $\Delta$ , define  $\Delta(\sigma)$  by induction as follows:

1.  $\Delta(\Lambda) = \Delta$
2.  $\Delta(\sigma.n) = \begin{cases} (\Delta(\sigma) \cup \{\psi_n\})^* & \text{if } \Delta(\sigma) \cup \{\psi_n\} \not\vdash \chi_n \\ \Delta(\sigma) & \text{otherwise} \end{cases}$

Here by  $(\Delta(\sigma) \cup \{\psi_n\})^*$  we mean the prime set of **formulas** which exists by ?? applied to the set  $\Delta(\sigma) \cup \{\psi_n\}$  and the **formula**  $\chi_n$ . Note that by this definition, if  $\Delta(\sigma) \cup \{\psi_n\} \not\vdash \chi_n$ , then  $\Delta(\sigma.n) \vdash \psi_n$  and  $\Delta(\sigma.n) \not\vdash \chi_n$ . Note also that  $\Delta(\sigma) \subseteq \Delta(\sigma.n)$  for any  $n$ . If  $\Delta$  is prime, then  $\Delta(\sigma)$  is prime for all  $\sigma$ .

int:sc:mod:  
defn:canonical-model **Definition sc.1.** Suppose  $\Delta$  is prime. Then the *canonical model*  $\mathfrak{M}(\Delta)$  for  $\Delta$  is defined by:

1.  $W = \mathbb{N}^*$ , the set of finite sequences of natural numbers.
2.  $R$  is the partial order according to which  $R\sigma\sigma'$  iff  $\sigma$  is an initial segment of  $\sigma'$  (i.e.,  $\sigma' = \sigma \frown \sigma''$  for some sequence  $\sigma''$ ).
3.  $V(p) = \{\sigma : p \in \Delta(\sigma)\}$ .

It is easy to verify that  $R$  is indeed a partial order. Also, the monotonicity condition on  $V$  is satisfied. Since  $\Delta(\sigma) \subseteq \Delta(\sigma.n)$  we get  $\Delta(\sigma) \subseteq \Delta(\sigma')$  whenever  $R\sigma\sigma'$  by induction on  $\sigma$ .

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## Bibliography