

int.1 Sentences

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Ok, now we have a (sketch of a) definition of satisfaction (“true in”) for **structures** and **formulas**. But it needs this additional bit—a **variable** assignment—and what we wanted is a definition of **sentences**. How do we get rid of assignments, and what are **sentences**?

You probably remember a discussion in your first introduction to formal logic about the relation between **variables** and quantifiers. A quantifier is always followed by a **variable**, and then in the part of the **sentence** to which that quantifier applies (its “scope”), we understand that the **variable** is “bound” by that quantifier. In **formulas** it was not required that every **variable** has a matching quantifier, and **variables** without matching quantifiers are “free” or “unbound.” We will take **sentences** to be all those **formulas** that have no free **variables**.

Again, the intuitive idea of when an occurrence of a **variable** in a **formula** φ is bound, which quantifier binds it, and when it is free, is not difficult to get. You may have learned a method for testing this, perhaps involving counting parentheses. We have to insist on a precise definition—and because we have defined **formulas** by induction, we can give a definition of the free and bound occurrences of a **variable** x in a **formula** φ also by induction. E.g., it might look like this for our simplified language:

1. If φ is atomic, all occurrences of x in it are free (that is, the occurrence of x in $P(x)$ is free).
2. If φ is of the form $\neg\psi$, then an occurrence of x in $\neg\psi$ is free iff the corresponding occurrence of x is free in ψ (that is, the free occurrences of variables in ψ are exactly the corresponding occurrences in $\neg\psi$).
3. If φ is of the form $(\psi \wedge \chi)$, then an occurrence of x in $(\psi \wedge \chi)$ is free iff the corresponding occurrence of x is free in ψ or in χ .
4. If φ is of the form $\exists x \psi$, then no occurrence of x in φ is free; if it is of the form $\exists y \psi$ where y is a different **variable** than x , then an occurrence of x in $\exists y \psi$ is free iff the corresponding occurrence of x is free in ψ .

Once we have a precise definition of free and bound occurrences of variables, we can simply say: a **sentence** is any **formula** without free occurrences of **variables**.

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Bibliography