

Chapter udf

The Iterative Conception

story.1 Extensionality

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The very first thing to say is that sets are individuated by their **elements**. More precisely:

Axiom (Extensionality). If sets A and B have the same **elements**, then A and B are the same set.

$$\forall A \forall B (\forall x (x \in A \leftrightarrow x \in B) \rightarrow A = B)$$

We assumed this throughout ???. But it bears repeating. The Axiom of Extensionality expresses the basic idea that a set is determined by its **elements**. (So sets might be contrasted with *concepts*, where precisely the same objects might fall under many different concepts.)

Why embrace this principle? Well, it is plausible to say that any denial of Extensionality is a decision to abandon anything which might even be called *set theory*. Set theory is no more nor less than the theory of extensional collections.

The real challenge in ???, though, is to lay down principles which tell us *which sets exist*. And it turns out that the only truly “obvious” answer to this question is provably wrong.

story.2 Russell’s Paradox (again)

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In ???, we worked with a naïve set theory. But according to a *very* naïve conception, sets are just the extensions of predicates. This naïve thought would mandate the following principle:

Naïve Comprehension. $\{x : \varphi(x)\}$ exists for any formula φ .

Tempting as this principle is, it is provably inconsistent. We saw this in ???, but the result is so important, and so straightforward, that it’s worth repeating. Verbatim.

Theorem story.1 (Russell’s Paradox). *There is no set $R = \{x : x \notin x\}$*

Proof. If $R = \{x : x \notin x\}$ exists, then $R \in R$ iff $R \notin R$, which is a contradiction. \square

Russell discovered this result in June 1901. (He did not, though, put the paradox in quite the form we just presented it, since he was considering Frege’s set theory, as outlined in *Grundgesetze*. We will return to this in [section story.6](#).) Russell wrote to Frege on June 16, 1902, explaining the inconsistency in Frege’s system. For the correspondence, and a bit of background, see [Heijenoort \(1967, pp. 124–8\)](#).

It is worth emphasising that this two-line proof is a result of *pure logic*. Granted, we implicitly used a (non-logical?) axiom, Extensionality, in our notation $\{x : x \notin x\}$; for $\{x : \varphi(x)\}$ is to be *the unique* (by Extensionality) set of the φ s, if one exists. But we can avoid even the hint of Extensionality, just by stating the result as follows: *there is no set whose members are exactly the non-self-membered sets*. And this has nothing much to do with sets. As Russell himself observed, exactly similar reasoning will lead you to conclude: *no man shaves exactly the men who do not shave themselves. Or: no pug sniffs exactly the pugs which don’t sniff themselves*. And so on. Schematically, the shape of the result is just:

$$\neg \exists x \forall z (Rzx \leftrightarrow \neg Rzz).$$

And that’s just a theorem (scheme) of first-order logic. Consequently, we can’t avoid Russell’s Paradox just by tinkering with our set theory; it arises before we even *get* to set theory. If we’re going to use (classical) first-order logic, we simply have to *accept* that there is no set $R = \{x : x \notin x\}$.

The upshot is this. If you want to accept Naïve Comprehension whilst *avoiding* inconsistency, you cannot just tinker with the *set theory*. Instead, you would have to overhaul your *logic*.

Of course, set theories with non-classical logics have been presented. But they are—to say the least—non-standard. The standard approach to Russell’s Paradox is to treat it as a straightforward non-existence proof, and then to try to learn how to live with it. That is the approach we will follow.

story.3 Predicative and Impredicative

The Russell set, R , was defined via $\{x : x \notin x\}$. Spelled out more fully, R would be the set which contains all and only those sets which are not non-self-membered. So in defining R , we quantify over the domain which would contain R (if it existed).

This is an *impredicative* definition. More generally, we might say that a definition is impredicative iff it quantifies over a domain which contains the object that is being defined.

In the wake of the paradoxes, Whitehead, Russell, Poincaré and Weyl rejected such impredicative definitions as “viciously circular”:

An analysis of the paradoxes to be avoided shows that they all result from a kind of vicious circle. The vicious circles in question arise from supposing that a collection of objects may contain members which can only be defined by means of the collection as a whole[. . . ¶]

The principle which enables us to avoid illegitimate totalities may be stated as follows: ‘Whatever involves *all* of a collection must not be one of the collection’; or, conversely: ‘If, provided a certain collection had a total, it would have members only definable in terms of that total, then the said collection has no total.’ We shall call this the ‘vicious-circle principle,’ because it enables us to avoid the vicious circles involved in the assumption of illegitimate totalities. (Whitehead and Russell, 1910, p. 37)

If we follow them in rejecting the *vicious-circle principle*, then we might attempt to replace the disastrous Naïve Comprehension Scheme (of [section story.2](#)) with something like this:

Predicative Comprehension. For every formula φ quantifying only over sets: the set' $\{x : \varphi(x)\}$ exists.

So long as sets' are not sets, no contradiction will ensue.

Unfortunately, Predicative Comprehension is not very *comprehensive*. After all, it introduces us to new entities, sets'. So we will have to consider formulas which quantify over sets'. If they always yield a set', then Russell's paradox will arise again, just by considering the set' of all non-self-membered sets'. So, pursuing the same thought, we must say that a formula quantifying over sets' yields a corresponding set''. And then we will need sets''', sets''', etc. To prevent a rash of primes, it will be easier to think of these as sets₀, sets₁, sets₂, sets₃, sets₄, And this would give us a way into the (simple) theory of types.

There are a few obvious objections against such a theory (though it is not obvious that they are *overwhelming* objections). In brief: the resulting theory is cumbersome to use; it is profligate in postulating different kinds of objects; and it is not clear, in the end, that impredicative definitions are even *all that bad*.

To bring out the last point, consider this remark from Ramsey:

we may refer to a man as the tallest in a group, thus identifying him by means of a totality of which he is himself a member without there being any vicious circle. (Ramsey, 1925)

Ramsey's point is that “the tallest man in the group” *is* an impredicative definition; but it is obviously perfectly kosher.

One might respond that, in this case, we could pick out the tallest person by *predicative* means. For example, maybe we could just point at the man in question. The objection against impredicative definitions, then, would clearly need

to be limited to entities which can *only* be picked out impredicatively. But even then, we would need to hear more, about why such “essential impredicativity” would be so bad.¹

Admittedly, impredicative definitions are extremely bad news, if we want our definitions to provide us with something like a recipe for *creating* an object. For, given an impredicative definition, one would genuinely be caught in a vicious circle: to create the impredicatively specified object, one would *first* need to create all the objects (including the impredicatively specified object), since the impredicatively specified object is specified in terms of all the objects; so one would need to create the impredicatively specified object before one had created it itself. But again, this is only a serious objection against “essentially impredicatively” specified sets, if we think of sets as things that we *create*. And we (probably) don’t.

As such—for better or worse—the approach which became common does not involve taking a hard line concerning (im)predicativity. Rather, it involves what is now regarded as the cumulative-iterative approach. In the end, this will allow us to stratify our sets into “stages”—a *bit* like the predicative approach stratifies entities into $\text{sets}_0, \text{sets}_1, \text{sets}_2, \dots$ —but we will not postulate any difference in kind between them.

story.4 The Cumulative-Iterative Approach

Here is a slightly fuller statement of how we will stratify sets into stages:

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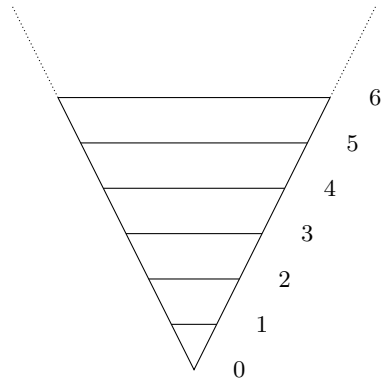
Sets are formed in *stages*. For each stage S , there are certain stages which are *before* S . At stage S , each collection consisting of sets formed at stages before S is formed into a set. There are no sets other than the sets which are formed at stages. (Shoenfield, 1977, p. 323)

This is a sketch of the *cumulative-iterative conception of set*. It will underpin the formal set theory that we present in ??.

Let’s explore this in a little more detail. As Shoenfield describes the process, at every stage, we form new sets from the sets which were available to us from earlier stages. So, on Shoenfield’s picture, at the initial stage, stage 0, there are no *earlier* stages, and so *a fortiori* there are no sets available to us from earlier stages.² So we form only one set: the set with no **elements** \emptyset . At stage 1, exactly one set is available to us from earlier stages, so only one new set is $\{\emptyset\}$. At stage 2, two sets are available to us from earlier stages, and we form two new sets $\{\{\emptyset\}\}$ and $\{\emptyset, \{\emptyset\}\}$. At stage 3, four sets are available to us from earlier stages, so we form twelve new sets. . . . As such, the cumulative-iterative picture of the sets will look a bit like this (with numbers indicating stages):

¹For more, see Linnebo (2010).

²Why should we assume that there *is* a first stage? See the footnote to *Stages-are-ordered* in ??.



So: why should we embrace this story?

One reason is that it is a nice, tractable story. Given the demise of the most obvious story, i.e., Naïve Comprehension, we are in want of something nice.

But the story is not *just* nice. We have a good reason to believe that any set theory based on this story will be *consistent*. Here is why.

Given the cumulative-iterative conception of set, we form sets at stages; and their **elements** must be objects which were available *already*. So, for any stage S , we can form the set

$$R_S = \{x : x \notin x \text{ and } x \text{ was available before } S\}$$

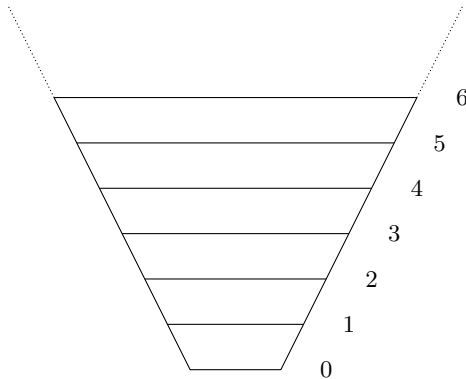
The reasoning involved in proving Russell’s Paradox will now establish that R_S itself is not available before stage S . And that’s not a contradiction. Moreover, if we embrace the cumulative-iterative conception of set, then we shouldn’t even have *expected* to be able to form the Russell set itself. For that would be the set of all non-self-membered sets that “will ever be available”. In short: the fact that we (provably) can’t form the Russell set isn’t *surprising*, given the cumulative-iterative story; it’s what we would *predict*.

story.5 Urelements or Not?

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In the next few chapters, we will try to extract axioms from the cumulative-iterative conception of set. But, before going any further, we need to say something more about *urelements*.

The picture of [section story.4](#) allowed us only to form new sets from old *sets*. However, we might want to allow that certain *non-sets*—cows, pigs, grains of sand, or whatever—can be **elements** of sets. In that case, we would start with certain basic elements, *urelements*, and then say that at each stage S we would form “all possible” sets consisting of urelements or sets formed at stages before S (in any combination). The resulting picture would look more like this:



So now we have a decision to take: *Should we allow urelements?*

Philosophically, it makes sense to include urelements in our theorising. The main reason for this is to make our set theory *applicable*. To illustrate the point, recall from ?? that we say that two sets A and B have the same size, i.e., $A \approx B$, iff there is a bijection between them. Now, if the cows in the field and the pigs in the sty both form sets, we can offer a set-theoretical treatment of the claim “there are as many cows as pigs”. But if we ban urelements, so that the cows and the pigs do *not* form sets, then that set-theoretical treatment will be unavailable. Indeed, we will have no straightforward ability to apply set theory to anything other than sets themselves. (For more reasons to include urelements, see [Potter 2004](#), pp. vi, 24, 50–1.)

Mathematically, however, it is quite rare to allow urelements. In part, this is because it is *very slightly* easier to formulate set theory without urelements. But, occasionally, one finds more interesting justifications for excluding urelement from set theory:

In accordance with the belief that set theory is the foundation of mathematics, we should be able to capture all of mathematics by just talking about sets, so our variable should not range over objects like cows and pigs. ([Kunen, 1980](#), p. 8)

So: a focus on applicability would suggest *including* urelements; a focus on a reductive foundational goal (reducing mathematics to pure set theory) might suggest *excluding* them. Mild laziness, too, points in the direction of excluding urelements.

We will follow the laziest path. Partly, though, there is a pedagogical justification. Our aim is to introduce you to the elements of set theory that you would need in order to get started on the philosophy of set theory. And most of that philosophical literature discusses set theories formulated *without* urelements. So this book will, perhaps, be of more use, if it hews fairly closely to that literature.

story.6 Appendix: Frege’s Basic Law V

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In [section story.2](#), we explained that Russell’s formulated his paradox as a problem for the system Frege outlined in his *Grundgesetze*. Frege’s system did not include a direct formulation of Naïve Comprehension. So, in this appendix, we will very briefly explain what Frege’s system *did* include, and how it relates to Naïve Comprehension and how it relates to Russell’s Paradox.

Frege’s system is *second-order*, and was designed to formulate the notion of an *extension of a concept*. Using notation inspired by Frege, we will write $\epsilon x F(x)$ for *the extension of the concept F*. This is a device which takes a *predicate*, “*F*”, and turns it into a (first-order) *term*, “ $\epsilon x F(x)$ ”. Using this device, Frege offered the following *definition* of membership:

$$a \in b =_{\text{df}} \exists G(b = \epsilon x G(x) \wedge Ga)$$

roughly: $a \in b$ iff a falls under a concept whose extension is b . (Note that the quantifier “ $\exists G$ ” is second-order.) Frege also maintained the following principle, known as *Basic Law V*:

$$\epsilon x F(x) = \epsilon x G(x) \leftrightarrow \forall x(Fx \leftrightarrow Gx)$$

roughly: concepts have identical extensions iff they are coextensive. (Again, both “*F*” and “*G*” are in predicate position.) Now a simple principle connects membership with property-satisfaction:

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Lemma story.2 (in *Grundgesetze*). $\forall F \forall a(a \in \epsilon x F(x) \leftrightarrow Fa)$

Proof. Fix F and a . Now $a \in \epsilon x F(x)$ iff $\exists G(\epsilon x F(x) = \epsilon x G(x) \wedge Ga)$ (by the definition of membership) iff $\exists G(\forall x(Fx \leftrightarrow Gx) \wedge Ga)$ (by Basic Law V) iff Fa (by elementary second-order logic). \square

And this yields Naïve Comprehension almost immediately:

Lemma story.3 (in *Grundgesetze*). $\forall F \exists s \forall a(a \in s \leftrightarrow Fa)$

Proof. Fix F ; now [Lemma story.2](#) yields $\forall a(a \in \epsilon x F(x) \leftrightarrow Fa)$; so $\exists s \forall a(a \in s \leftrightarrow Fa)$ by existential generalisation. The result follows since F was arbitrary. \square

Russell’s Paradox follows by taking F as given by $\forall x(Fx \leftrightarrow x \notin x)$.

Photo Credits

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