

choice.1 Comparability and Hartogs' Lemma

sth:choice:hartogs:
sec That's the plus side. Here's the minus side. Without Choice, things get *messy*. To see why, here is a nice result due to [Hartogs \(1915\)](#):

sth:choice:hartogs:
HartogsLemma **Lemma choice.1** (in **ZF**). *For any set A , there is an ordinal α such that $\alpha \not\preceq A$*

Proof. If $B \subseteq A$ and $R \subseteq B^2$, then $\langle B, R \rangle \subseteq V_{\text{rank}(A)+4}$ by [??](#). So, using Separation, consider:

$$C = \{\langle B, R \rangle \in V_{\text{rank}(A)+5} : B \subseteq A \text{ and } \langle B, R \rangle \text{ is a well-ordering}\}$$

Using Replacement and [??](#), form the set:

$$\alpha = \{\text{ord}(B, R) : \langle B, R \rangle \in C\}.$$

By [??](#), α is an ordinal, since it is a transitive set of ordinals. After all, if $\gamma \in \beta \in \alpha$, then $\beta = \text{ord}(B, R)$ for some $B \subseteq R$, whereupon $\gamma = \text{ord}(B_b, R_b)$ for some $b \in B$ by [??](#), so that $\gamma \in \alpha$.

For reductio, suppose there is an [injection](#) $f: \alpha \rightarrow A$. Then, where:

$$\begin{aligned} B &= \text{ran}(f) \\ R &= \{\langle f(\alpha), f(\beta) \rangle \in A \times A : \alpha \in \beta\}. \end{aligned}$$

Clearly $\alpha = \text{ord}(B, R)$ and $\langle B, R \rangle \in C$. So $\alpha \in \alpha$, which is a contradiction. \square

This entails a deep result:

Theorem choice.2 (in **ZF**). *The following claims are equivalent:*

- sth:choice:hartogs:
equivwo 1. *The Axiom of Well-Ordering*
- sth:choice:hartogs:
equivcompare 2. *Either $A \preceq B$ or $B \preceq A$, for any sets A and B*

Proof. (1) \Rightarrow (2). Fix A and B . Invoking (1), there are well-orderings $\langle A, R \rangle$ and $\langle B, S \rangle$. Invoking [??](#), let $f: \alpha \rightarrow \langle A, R \rangle$ and $g: \beta \rightarrow \langle B, S \rangle$ be isomorphisms. By Trichotomy, either $\alpha \in \beta$ or $\alpha = \beta$ or $\beta \in \alpha$. In the first two cases $\alpha \subseteq \beta$, so $g \circ f^{-1}: A \rightarrow B$ is an [injection](#), and hence $A \preceq B$. Similarly if $\beta \in \alpha$ then $B \preceq A$.

(1) \Rightarrow (2). Fix A ; by [Lemma choice.1](#) there is some ordinal β such that $\beta \not\preceq A$. Invoking (2), we have $A \preceq \beta$. So there is some [injection](#) $f: A \rightarrow \beta$, and we can use this injection to well-order the elements of A , by defining an order $\{\langle a, b \rangle \in A \times A : f(a) \in f(b)\}$. \square

As an immediate consequence: if Well-Ordering fails, then some sets are *literally incomparable* with regard to their size. So, if Well-Ordering fails, then transfinite cardinal arithmetic will be messy. For example, we will have to abandon the idea that if A and B are infinite then $A \sqcup B \approx A \times B \approx M$, where

M is the larger of A and B (see ??). The problem is simple: if we cannot *compare* the size of A and B , then it is nonsensical to ask which is larger.

Photo Credits

Bibliography

Hartogs, Friedrich. 1915. Über das Problem der Wohlordnung. *Mathematische Annalen* 76: 438–43.