

fil.1 Filtrations

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sec

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def:filtration **Definition fil.1.** Let Γ be closed under subformulas and $\mathfrak{M} = \langle W, R, V \rangle$. A *filtration of \mathfrak{M} through Γ* is any model $\mathfrak{M}^* = \langle W^*, R^*, V^* \rangle$, where:

1. $W^* = \{[w] : w \in W\}$;
2. For any $u, v \in W$:
 - a) If Ruv then $R^*[u][v]$;
 - b) If $R^*[u][v]$ then for any $\Box\varphi \in \Gamma$, if $\mathfrak{M}, u \models \Box\varphi$ then $\mathfrak{M}, v \models \varphi$;
 - c) If $R^*[u][v]$ then for any $\Diamond\varphi \in \Gamma$, if $\mathfrak{M}, v \models \varphi$ then $\mathfrak{M}, u \models \Diamond\varphi$.
3. $V^*(p) = \{[u] : u \in V(p)\}$.

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def:filtration-R
mod:fil:fil:
def:filtration-R1
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def:filtration-R2
mod:fil:fil:
def:filtration-R3

mod:fil:fil:
thm:filtrations **Theorem fil.2.** *If \mathfrak{M}^* is a filtration of \mathfrak{M} through Γ , then for every $\varphi \in \Gamma$ and $w \in W$, we have $\mathfrak{M}, w \models \varphi$ if and only if $\mathfrak{M}^*, [w] \models \varphi$.*

Proof. By induction on φ , using the fact that Γ is closed under subformulas. For φ atomic, say p : the left-to-right direction is immediate, as $\mathfrak{M}, w \models p$ only if $w \in V(p)$, which implies $[w] \in V^*(p)$, i.e., $\mathfrak{M}^*, [w] \models p$. Conversely, suppose $\mathfrak{M}^*, [w] \models p$, i.e., $[w] \in V^*(p)$; then $w \equiv w' \in V(p)$, and since $p \in \Gamma$, also $w \in V(p)$, so that $\mathfrak{M}, w \models p$. The cases for the Boolean connectives follow immediately from the inductive hypothesis and closure of Γ under subformulas.

So we do the case for $\Box\varphi \in \Gamma$. Suppose $\mathfrak{M}, u \models \Box\varphi$; to show that $\mathfrak{M}^*, [u] \models \varphi$, let v be such that $R^*[u][v]$. From **Definition fil.1(2b)**, we have that $\mathfrak{M}, v \models \varphi$, and by inductive hypothesis $\mathfrak{M}^*, [v] \models \varphi$. Since v was arbitrary, $\mathfrak{M}^*, [u] \models \Box\varphi$ follows. Conversely, suppose $\mathfrak{M}^*, [u] \models \Box\varphi$ and let v be arbitrary such that Ruv . From **Definition fil.1(2a)**, we have $R^*[u][v]$, so that $\mathfrak{M}^*, [v] \models \varphi$; by inductive hypothesis $\mathfrak{M}, v \models \varphi$, and since v was arbitrary, $\mathfrak{M}, u \models \Box\varphi$. \square

Corollary fil.3. *Let Γ be closed under subformulas. Then:*

1. *If \mathfrak{M}^* is a filtration of \mathfrak{M} through Γ then for any $\varphi \in \Gamma$: $\mathfrak{M} \models \varphi$ if and only if $\mathfrak{M}^* \models \varphi$.*
2. *If \mathcal{C} is a class of models and $\Gamma(\mathcal{C})$ is the class of Γ -filtrations of models in \mathcal{C} , then any formula $\varphi \in \Gamma$ is valid in \mathcal{C} if and only if it is valid in $\Gamma(\mathcal{C})$.*

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Bibliography