

## com.1 Lindenbaum's Lemma

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thm:lindenbaum **Theorem com.1** (Lindenbaum's Lemma). *If  $\Gamma$  is  $\Sigma$ -consistent then there is a complete  $\Sigma$ -consistent set  $\Delta$  extending  $\Gamma$ .*

*Proof.* Let  $\varphi_0, \varphi_1, \dots$  be an exhaustive listing of all formulas of the language (repetitions are allowed). For instance, start by listing  $p_0$ , and at each stage  $n$  list the finitely many formulas of length  $n$  using only variables among  $p_0, \dots, p_n$ . We define sets of formulas  $\Delta_n$  by induction on  $n$ , and we then set  $\Delta = \bigcup_n \Delta_n$ . We first put  $\Delta_0 = \Gamma$ , then supposing that  $\Delta_n$  has been defined:

$$\Delta_{n+1} = \begin{cases} \Delta_n \cup \{\varphi_n\}, & \text{if } \Delta_n \cup \{\varphi_n\} \text{ is consistent;} \\ \Delta_n \cup \{\neg\varphi_n\}, & \text{otherwise.} \end{cases}$$

If we now let  $\Delta = \bigcup_n \Delta_n$ , we can show the following:

1. For each  $n$ ,  $\Delta_n \subseteq \Delta$  (immediate from the definition).
2.  $\Gamma \subseteq \Delta$  (from (a)).
3. If  $n \leq m$  then  $\Delta_n \subseteq \Delta_m$  (by induction on  $m - n$ ).
4.  $\Delta$  is maximal (by construction).
5. For each  $m$ ,  $\Delta_m$  is consistent (by induction on  $m$ , using ????)
6. If  $\Delta' \subseteq \Delta$  is finite, then there is  $m$  such that  $\Delta' \subseteq \Delta_m$ .
7.  $\Delta$  is consistent.

It follows that  $\Delta$  is a complete  $\Sigma$ -consistent set extending  $\Gamma$ . □

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## Bibliography