

com.1 Determination and Completeness for \mathbf{K}

mod:com:cmk:
sec We are now prepared to use the canonical model to establish completeness. Completeness follows from the fact that the formulas true in the canonical for Σ are exactly the Σ -derivable ones. Models with this property are said to *determine* Σ .

Definition com.1. A model \mathfrak{M} *determines* a normal modal logic Σ precisely when $\mathfrak{M} \Vdash \varphi$ if and only if $\Sigma \vdash \varphi$, for all formulas φ .

mod:com:cmk:
thm:determination **Theorem com.2** (Determination). $\mathfrak{M}^\Sigma \Vdash \varphi$ if and only if $\Sigma \vdash \varphi$.

Proof. If $\mathfrak{M}^\Sigma \Vdash \varphi$, then for every complete Σ -consistent Δ , we have $\mathfrak{M}^\Sigma, \Delta \Vdash \varphi$. Hence, by the Truth Lemma, $\varphi \in \Delta$ for every complete Σ -consistent Δ , whence by ?? (with $\Gamma = \emptyset$), $\Sigma \vdash \varphi$.

Conversely, if $\Sigma \vdash \varphi$ then by ?????, every complete Σ -consistent Δ contains φ , and hence by the Truth Lemma, $\mathfrak{M}^\Sigma, \Delta \Vdash \varphi$ for every $\Delta \in W^\Sigma$, i.e., $\mathfrak{M}^\Sigma \Vdash \varphi$. \square

Since the canonical model for \mathbf{K} determines \mathbf{K} , we immediately have completeness of \mathbf{K} as a corollary:

mod:com:cmk:
cor:Kcomplete **Corollary com.3.** *The basic modal logic \mathbf{K} is complete with respect to the class of all models, i.e., if $\models \varphi$ then $\mathbf{K} \vdash \varphi$.*

Proof. Contrapositively, if $\mathbf{K} \not\vdash \varphi$ then by Determination $\mathfrak{M}^\mathbf{K} \not\vdash \varphi$ and hence φ is not valid. \square

For the general case of completeness of a system Σ with respect to a class of models, e.g., of $\mathbf{KTB4}$ with respect to the class of reflexive, symmetric, transitive models, determination alone is not enough. We must also show that the canonical model for the system Σ is a member of the class, which does not follow obviously from the canonical model construction—nor is it always true!

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Bibliography