

## mar.1 Standard Models of Arithmetic

mod:mar:stm:  
sec

The language of arithmetic  $\mathcal{L}_A$  is obviously intended to be about numbers, specifically, about natural numbers. So, “the” standard model  $\mathfrak{N}$  is special: it is the model we want to talk about. But in logic, we are often just interested in structural properties, and any two **structures** that are isomorphic share those. So we can be a bit more liberal, and consider any **structure** that is isomorphic to  $\mathfrak{N}$  “standard.”

**Definition mar.1.** A **structure** for  $\mathcal{L}_A$  is *standard* if it is isomorphic to  $\mathfrak{N}$ .

mod:mar:stm:  
prop:standard-domain

**Proposition mar.2.** *If a **structure**  $\mathfrak{M}$  is standard, then its domain is the set of values of the standard numerals, i.e.,*

$$|\mathfrak{M}| = \{\text{Val}^{\mathfrak{M}}(\bar{n}) : n \in \mathbb{N}\}$$

*Proof.* Clearly, every  $\text{Val}^{\mathfrak{M}}(\bar{n}) \in |\mathfrak{M}|$ . We just have to show that every  $x \in |\mathfrak{M}|$  is equal to  $\text{Val}^{\mathfrak{M}}(\bar{n})$  for some  $n$ . Since  $\mathfrak{M}$  is standard, it is isomorphic to  $\mathfrak{N}$ . Suppose  $g: \mathbb{N} \rightarrow |\mathfrak{M}|$  is an isomorphism. Then  $g(n) = g(\text{Val}^{\mathfrak{N}}(\bar{n})) = \text{Val}^{\mathfrak{M}}(\bar{n})$ . But for every  $x \in |\mathfrak{M}|$ , there is an  $n \in \mathbb{N}$  such that  $g(n) = x$ , since  $g$  is **surjective**.  $\square$

If a structure  $\mathfrak{M}$  for  $\mathcal{L}_A$  is standard, the elements of its **domain** can all be named by the standard numerals  $\bar{0}, \bar{1}, \bar{2}, \dots$ , i.e., the terms  $o, o', o'', \dots$ . Of course, this does not mean that the **elements** of  $|\mathfrak{M}|$  *are* the numbers, just that we can pick them out the same way we can pick out the numbers in  $|\mathfrak{N}|$ . explanation

**Problem mar.1.** Show that the converse of **Proposition mar.2** is false, i.e., give an example of a **structure**  $\mathfrak{M}$  with  $|\mathfrak{M}| = \{\text{Val}^{\mathfrak{M}}(\bar{n}) : n \in \mathbb{N}\}$  that is not isomorphic to  $\mathfrak{N}$ .

mod:mar:stm:  
prop:thq-standard

**Proposition mar.3.** *If  $\mathfrak{M} \models \mathbf{Q}$ , and  $|\mathfrak{M}| = \{\text{Val}^{\mathfrak{M}}(\bar{n}) : n \in \mathbb{N}\}$ , then  $\mathfrak{M}$  is standard.*

*Proof.* We have to show that  $\mathfrak{M}$  is isomorphic to  $\mathfrak{N}$ . Consider the function  $g: \mathbb{N} \rightarrow |\mathfrak{M}|$  defined by  $g(n) = \text{Val}^{\mathfrak{M}}(\bar{n})$ . By the hypothesis,  $g$  is **surjective**. It is also **injective**:  $\mathbf{Q} \vdash \bar{n} \neq \bar{m}$  whenever  $n \neq m$ . Thus, since  $\mathfrak{M} \models \mathbf{Q}$ ,  $\mathfrak{M} \models \bar{n} \neq \bar{m}$ , whenever  $n \neq m$ . Thus, if  $n \neq m$ , then  $\text{Val}^{\mathfrak{M}}(\bar{n}) \neq \text{Val}^{\mathfrak{M}}(\bar{m})$ , i.e.,  $g(n) \neq g(m)$ .

We also have to verify that  $g$  is an isomorphism.

1. We have  $g(o^{\mathfrak{N}}) = g(0)$  since,  $o^{\mathfrak{N}} = 0$ . By definition of  $g$ ,  $g(0) = \text{Val}^{\mathfrak{M}}(\bar{0})$ . But  $\bar{0}$  is just  $o$ , and the value of a term which happens to be a **constant symbol** is given by what the **structure** assigns to that **constant symbol**, i.e.,  $\text{Val}^{\mathfrak{M}}(o) = o^{\mathfrak{M}}$ . So we have  $g(o^{\mathfrak{N}}) = o^{\mathfrak{M}}$  as required.
2.  $g(\iota^{\mathfrak{N}}(n)) = g(n+1)$ , since  $\iota$  in  $\mathfrak{N}$  is the successor function on  $\mathbb{N}$ . Then,  $g(n+1) = \text{Val}^{\mathfrak{M}}(\overline{n+1})$  by definition of  $g$ . But  $\overline{n+1}$  is the same term as  $\bar{n}'$ , so  $\text{Val}^{\mathfrak{M}}(\overline{n+1}) = \text{Val}^{\mathfrak{M}}(\bar{n}')$ . By the definition of the value function, this is  $\iota^{\mathfrak{M}}(\text{Val}^{\mathfrak{M}}(\bar{n}))$ . Since  $\text{Val}^{\mathfrak{M}}(\bar{n}) = g(n)$  we get  $g(\iota^{\mathfrak{N}}(n)) = \iota^{\mathfrak{M}}(g(n))$ .

3.  $g(+^{\mathfrak{N}}(n, m)) = g(n + m)$ , since  $+$  in  $\mathfrak{N}$  is the addition function on  $\mathbb{N}$ . Then,  $g(n + m) = \text{Val}^{\mathfrak{M}}(\overline{n + m})$  by definition of  $g$ . But  $\mathbf{Q} \vdash \overline{n + m} = (\overline{n} + \overline{m})$ , so  $\text{Val}^{\mathfrak{M}}(\overline{n + m}) = \text{Val}^{\mathfrak{M}}(\overline{n} + \overline{m})$ . By the definition of the value function, this is  $= +^{\mathfrak{M}}(\text{Val}^{\mathfrak{M}}(\overline{n}), \text{Val}^{\mathfrak{M}}(\overline{m}))$ . Since  $\text{Val}^{\mathfrak{M}}(\overline{n}) = g(n)$  and  $\text{Val}^{\mathfrak{M}}(\overline{m}) = g(m)$ , we get  $g(+^{\mathfrak{N}}(n, m)) = +^{\mathfrak{M}}(g(n), g(m))$ .
4.  $g(\times^{\mathfrak{N}}(n, m)) = \times^{\mathfrak{M}}(g(n), g(m))$ : Exercise.
5.  $\langle n, m \rangle \in <^{\mathfrak{N}}$  iff  $n < m$ . If  $n < m$ , then  $\mathbf{Q} \vdash \overline{n} < \overline{m}$ , and also  $\mathfrak{M} \models \overline{n} < \overline{m}$ . Thus  $\langle \text{Val}^{\mathfrak{M}}(\overline{n}), \text{Val}^{\mathfrak{M}}(\overline{m}) \rangle \in <^{\mathfrak{M}}$ , i.e.,  $\langle g(n), g(m) \rangle \in <^{\mathfrak{M}}$ . If  $n \not< m$ , then  $\mathbf{Q} \vdash \neg \overline{n} < \overline{m}$ , and consequently  $\mathfrak{M} \not\models \overline{n} < \overline{m}$ . Thus, as before,  $\langle g(n), g(m) \rangle \notin <^{\mathfrak{M}}$ . Together, we get:  $\langle n, m \rangle \in <^{\mathfrak{N}}$  iff  $\langle g(n), g(m) \rangle \in <^{\mathfrak{M}}$ .  $\square$

explanation

The function  $g$  is the most obvious way of defining a mapping from  $\mathbb{N}$  to the domain of any other structure  $\mathfrak{M}$  for  $\mathcal{L}_A$ , since every such  $\mathfrak{M}$  contains elements named by  $\overline{0}, \overline{1}, \overline{2}$ , etc. So it isn't surprising that if  $\mathfrak{M}$  makes at least some basic statements about the  $\overline{n}$ 's true in the same way that  $\mathfrak{N}$  does, and  $g$  is also bijective, then  $g$  will turn into an isomorphism. In fact, if  $|\mathfrak{M}|$  contains no elements other than what the  $\overline{n}$ 's name, it's the only one.

**Proposition mar.4.** *If  $\mathfrak{M}$  is standard, then  $g$  from the proof of Proposition mar.3 is the only isomorphism from  $\mathfrak{N}$  to  $\mathfrak{M}$ .*

mod:mar:stm:  
prop:thq-unique-iso

*Proof.* Suppose  $h: \mathbb{N} \rightarrow |\mathfrak{M}|$  is an isomorphism between  $\mathfrak{N}$  and  $\mathfrak{M}$ . We show that  $g = h$  by induction on  $n$ . If  $n = 0$ , then  $g(0) = \mathfrak{o}^{\mathfrak{M}}$  by definition of  $g$ . But since  $h$  is an isomorphism,  $h(0) = h(\mathfrak{o}^{\mathfrak{N}}) = \mathfrak{o}^{\mathfrak{M}}$ , so  $g(0) = h(0)$ .

Now consider the case for  $n + 1$ . We have

$$\begin{aligned}
g(n + 1) &= \text{Val}^{\mathfrak{M}}(\overline{n + 1}) \text{ by definition of } g \\
&= \text{Val}^{\mathfrak{M}}(\overline{n'}) \text{ since } \overline{n + 1} \equiv \overline{n'} \\
&= \mathfrak{r}^{\mathfrak{M}}(\text{Val}^{\mathfrak{M}}(\overline{n})) \text{ by definition of } \text{Val}^{\mathfrak{M}}(t') \\
&= \mathfrak{r}^{\mathfrak{M}}(g(n)) \text{ by definition of } g \\
&= \mathfrak{r}^{\mathfrak{M}}(h(n)) \text{ by induction hypothesis} \\
&= h(\mathfrak{r}^{\mathfrak{N}}(n)) \text{ since } h \text{ is an isomorphism} \\
&= h(n + 1) \quad \square
\end{aligned}$$

explanation

For any denumerable set  $M$ , there's a bijection between  $\mathbb{N}$  and  $M$ , so every such set  $M$  is potentially the domain of a standard model  $\mathfrak{M}$ . In fact, once you pick an object  $z \in M$  and a suitable function  $s$  as  $\mathfrak{o}^{\mathfrak{M}}$  and  $\mathfrak{r}^{\mathfrak{M}}$ , the interpretations of  $+$ ,  $\times$ , and  $<$  is already fixed. Only functions  $s: M \rightarrow M \setminus \{z\}$  that are both injective and surjective are suitable in a standard model as  $\mathfrak{r}^{\mathfrak{M}}$ . The range of  $s$  cannot contain  $z$ , since otherwise  $\forall x \mathfrak{o} \neq x'$  would be false. That sentence is true in  $\mathfrak{N}$ , and so  $\mathfrak{M}$  also has to make it true. The function  $s$  has to be injective, since the successor function  $\mathfrak{r}^{\mathfrak{N}}$  in  $\mathfrak{N}$  is, and that  $\mathfrak{r}^{\mathfrak{M}}$  is injective

is expressed by a sentence true in  $\mathfrak{N}$ . It has to be surjective because otherwise there would be some  $x \in M \setminus \{z\}$  not in the domain of  $s$ , i.e., the sentence  $\forall x (x = 0 \vee \exists y y' = x)$  would be false in  $\mathfrak{M}$ —but it is true in  $\mathfrak{N}$ .

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## Bibliography