

mar.1 Standard Models of Arithmetic

mod:mar:stm:
sec

The language of arithmetic \mathcal{L}_A is obviously intended to be about numbers, specifically, about natural numbers. So, “the” standard model \mathfrak{N} is special: it is the model we want to talk about. But in logic, we are often just interested in structural properties, and any two **structures** that are isomorphic share those. So we can be a bit more liberal, and consider any **structure** that is isomorphic to \mathfrak{N} “standard.”

Definition mar.1. A **structure** for \mathcal{L}_A is *standard* if it is isomorphic to \mathfrak{N} .

mod:mar:stm:
prop:standard-domain

Proposition mar.2. *If a **structure** \mathfrak{M} is standard, then its domain is the set of values of the standard numerals, i.e.,*

$$|\mathfrak{M}| = \{\text{Val}^{\mathfrak{M}}(\bar{n}) : n \in \mathbb{N}\}$$

Proof. Clearly, every $\text{Val}^{\mathfrak{M}}(\bar{n}) \in |\mathfrak{M}|$. We just have to show that every $x \in |\mathfrak{M}|$ is equal to $\text{Val}^{\mathfrak{M}}(\bar{n})$ for some n . Since \mathfrak{M} is standard, it is isomorphic to \mathfrak{N} . Suppose $g: \mathbb{N} \rightarrow |\mathfrak{M}|$ is an isomorphism. Then $g(n) = g(\text{Val}^{\mathfrak{N}}(\bar{n})) = \text{Val}^{\mathfrak{M}}(\bar{n})$. But for every $x \in |\mathfrak{M}|$, there is an $n \in \mathbb{N}$ such that $g(n) = x$, since g is **surjective**. \square

If a structure \mathfrak{M} for \mathcal{L}_A is standard, the elements of its **domain** can all be named by the standard numerals $\bar{0}, \bar{1}, \bar{2}, \dots$, i.e., the terms o, o', o'' , etc. Of course, this does not mean that the **elements** of $|\mathfrak{M}|$ *are* the numbers, just that we can pick them out the same way we can pick out the numbers in $|\mathfrak{N}|$. explanation

Problem mar.1. Show that the converse of **Proposition mar.2** is false, i.e., give an example of a **structure** \mathfrak{M} with $|\mathfrak{M}| = \{\text{Val}^{\mathfrak{M}}(\bar{n}) : n \in \mathbb{N}\}$ that is not isomorphic to \mathfrak{N} .

mod:mar:stm:
prop:thq-standard

Proposition mar.3. *If $\mathfrak{M} \models \mathbf{Q}$, and $|\mathfrak{M}| = \{\text{Val}^{\mathfrak{M}}(\bar{n}) : n \in \mathbb{N}\}$, then \mathfrak{M} is standard.*

Proof. We have to show that \mathfrak{M} is isomorphic to \mathfrak{N} . Consider the function $g: \mathbb{N} \rightarrow |\mathfrak{M}|$ defined by $g(n) = \text{Val}^{\mathfrak{M}}(\bar{n})$. By the hypothesis, g is **surjective**. It is also **injective**: $\mathbf{Q} \vdash \bar{n} \neq \bar{m}$ whenever $n \neq m$. Thus, since $\mathfrak{M} \models \mathbf{Q}$, $\mathfrak{M} \models \bar{n} \neq \bar{m}$, whenever $n \neq m$. Thus, if $n \neq m$, then $\text{Val}^{\mathfrak{M}}(\bar{n}) \neq \text{Val}^{\mathfrak{M}}(\bar{m})$, i.e., $g(n) \neq g(m)$.

We also have to verify that g is an isomorphism.

1. We have $g(o^{\mathfrak{N}}) = g(0)$ since, $o^{\mathfrak{N}} = 0$. By definition of g , $g(0) = \text{Val}^{\mathfrak{M}}(\bar{0})$. But $\bar{0}$ is just o , and the value of a term which happens to be a **constant symbol** is given by what the **structure** assigns to that **constant symbol**, i.e., $\text{Val}^{\mathfrak{M}}(o) = o^{\mathfrak{M}}$. So we have $g(o^{\mathfrak{N}}) = o^{\mathfrak{M}}$ as required.
2. $g(\iota^{\mathfrak{N}}(n)) = g(n+1)$, since ι in \mathfrak{N} is the successor function on \mathbb{N} . Then, $g(n+1) = \text{Val}^{\mathfrak{M}}(\overline{n+1})$ by definition of g . But $\overline{n+1}$ is the same term as \bar{n}' , so $\text{Val}^{\mathfrak{M}}(\overline{n+1}) = \text{Val}^{\mathfrak{M}}(\bar{n}')$. By the definition of the value function, this is $\iota^{\mathfrak{M}}(\text{Val}^{\mathfrak{M}}(\bar{n}))$. Since $\text{Val}^{\mathfrak{M}}(\bar{n}) = g(n)$ we get $g(\iota^{\mathfrak{N}}(n)) = \iota^{\mathfrak{M}}(g(n))$.

3. $g(+^{\mathfrak{N}}(n, m)) = g(n + m)$, since $+$ in \mathfrak{N} is the addition function on \mathbb{N} . Then, $g(n + m) = \text{Val}^{\mathfrak{M}}(\overline{n + m})$ by definition of g . But $\mathbf{Q} \vdash \overline{n + m} = (\overline{n} + \overline{m})$, so $\text{Val}^{\mathfrak{M}}(\overline{n + m}) = \text{Val}^{\mathfrak{M}}(\overline{n} + \overline{m})$. By the definition of the value function, this is $= +^{\mathfrak{M}}(\text{Val}^{\mathfrak{M}}(\overline{n}), \text{Val}^{\mathfrak{M}}(\overline{m}))$. Since $\text{Val}^{\mathfrak{M}}(\overline{n}) = g(n)$ and $\text{Val}^{\mathfrak{M}}(\overline{m}) = g(m)$, we get $g(+^{\mathfrak{N}}(n, m)) = +^{\mathfrak{M}}(g(n), g(m))$.
4. $g(\times^{\mathfrak{N}}(n, m)) = \times^{\mathfrak{M}}(g(n), g(m))$: Exercise.
5. $\langle n, m \rangle \in <^{\mathfrak{N}}$ iff $n < m$. If $n < m$, then $\mathbf{Q} \vdash \overline{n} < \overline{m}$, and also $\mathfrak{M} \models \overline{n} < \overline{m}$. Thus $\langle \text{Val}^{\mathfrak{M}}(\overline{n}), \text{Val}^{\mathfrak{M}}(\overline{m}) \rangle \in <^{\mathfrak{M}}$, i.e., $\langle g(n), g(m) \rangle \in <^{\mathfrak{M}}$. If $n \not< m$, then $\mathbf{Q} \vdash \neg \overline{n} < \overline{m}$, and consequently $\mathfrak{M} \not\models \overline{n} < \overline{m}$. Thus, as before, $\langle g(n), g(m) \rangle \notin <^{\mathfrak{M}}$. Together, we get: $\langle n, m \rangle \in <^{\mathfrak{N}}$ iff $\langle g(n), g(m) \rangle \in <^{\mathfrak{M}}$. \square

explanation

The function g is the most obvious way of defining a mapping from \mathbb{N} to the domain of any other structure \mathfrak{M} for \mathcal{L}_A , since every such \mathfrak{M} contains elements named by $\overline{0}, \overline{1}, \overline{2}$, etc. So it isn't surprising that if \mathfrak{M} makes at least some basic statements about the \overline{n} 's true in the same way that \mathfrak{N} does, and g is also bijective, then g will turn into an isomorphism. In fact, if $|\mathfrak{M}|$ contains no elements other than what the \overline{n} 's name, it's the only one.

Proposition mar.4. *If \mathfrak{M} is standard, then g from the proof of Proposition mar.3 is the only isomorphism from \mathfrak{N} to \mathfrak{M} .*

mod:mar:stm:
prop:thq-unique-iso

Proof. Suppose $h: \mathbb{N} \rightarrow |\mathfrak{M}|$ is an isomorphism between \mathfrak{N} and \mathfrak{M} . We show that $g = h$ by induction on n . If $n = 0$, then $g(0) = \mathfrak{o}^{\mathfrak{M}}$ by definition of g . But since h is an isomorphism, $h(0) = h(\mathfrak{o}^{\mathfrak{N}}) = \mathfrak{o}^{\mathfrak{M}}$, so $g(0) = h(0)$.

Now consider the case for $n + 1$. We have

$$\begin{aligned}
g(n + 1) &= \text{Val}^{\mathfrak{M}}(\overline{n + 1}) \text{ by definition of } g \\
&= \text{Val}^{\mathfrak{M}}(\overline{n'}) \text{ since } \overline{n + 1} \equiv \overline{n'} \\
&= \iota^{\mathfrak{M}}(\text{Val}^{\mathfrak{M}}(\overline{n})) \text{ by definition of } \text{Val}^{\mathfrak{M}}(t') \\
&= \iota^{\mathfrak{M}}(g(n)) \text{ by definition of } g \\
&= \iota^{\mathfrak{M}}(h(n)) \text{ by induction hypothesis} \\
&= h(\iota^{\mathfrak{N}}(n)) \text{ since } h \text{ is an isomorphism} \\
&= h(n + 1) \quad \square
\end{aligned}$$

explanation

For any denumerable set M , there's a bijection between \mathbb{N} and M , so every such set M is potentially the domain of a standard model \mathfrak{M} . In fact, once you pick an object $z \in M$ and a suitable function s as $\mathfrak{o}^{\mathfrak{M}}$ and $\iota^{\mathfrak{M}}$, the interpretations of $+$, \times , and $<$ is already fixed. Only functions $s: M \rightarrow M \setminus \{z\}$ that are both injective and surjective are suitable in a standard model as $\iota^{\mathfrak{M}}$. The range of s cannot contain z , since otherwise $\forall x \mathfrak{o} \neq x'$ would be false. That sentence is true in \mathfrak{N} , and so \mathfrak{M} also has to make it true. The function s has to be injective, since the successor function $\iota^{\mathfrak{N}}$ in \mathfrak{N} is, and that $\iota^{\mathfrak{N}}$ is injective

is expressed by a sentence true in \mathfrak{N} . It has to be surjective because otherwise there would be some $x \in M \setminus \{z\}$ not in the domain of s , i.e., the sentence $\forall x (x = 0 \vee \exists y y' = x)$ would be false in \mathfrak{M} —but it is true in \mathfrak{N} .

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Bibliography