

## bas.1 Reducts and Expansions

Often it is useful or necessary to compare languages which have symbols in common, as well as **structures** for these languages. The most common case is when all the symbols in a **language**  $\mathcal{L}$  are also part of a **language**  $\mathcal{L}'$ , i.e.,  $\mathcal{L} \subseteq \mathcal{L}'$ . An  $\mathcal{L}$ -**structure**  $\mathfrak{M}$  can then always be expanded to an  $\mathcal{L}'$ -**structure** by adding interpretations of the additional symbols while leaving the interpretations of the common symbols the same. On the other hand, from an  $\mathcal{L}'$ -**structure**  $\mathfrak{M}'$  we can obtain an  $\mathcal{L}$ -**structure** simply by “forgetting” the interpretations of the symbols that do not occur in  $\mathcal{L}$ .

mod:bas:red: defn:reduct **Definition bas.1.** Suppose  $\mathcal{L} \subseteq \mathcal{L}'$ ,  $\mathfrak{M}$  is an  $\mathcal{L}$ -**structure** and  $\mathfrak{M}'$  is an  $\mathcal{L}'$ -**structure**.  $\mathfrak{M}$  is the *reduct* of  $\mathfrak{M}'$  to  $\mathcal{L}$ , and  $\mathfrak{M}'$  is an *expansion* of  $\mathfrak{M}$  to  $\mathcal{L}'$  iff

1.  $|\mathfrak{M}| = |\mathfrak{M}'|$
2. For every **constant symbol**  $c \in \mathcal{L}$ ,  $c^{\mathfrak{M}} = c^{\mathfrak{M}'}$ .
3. For every **function symbol**  $f \in \mathcal{L}$ ,  $f^{\mathfrak{M}} = f^{\mathfrak{M}'}$ .
4. For every **predicate symbol**  $P \in \mathcal{L}$ ,  $P^{\mathfrak{M}} = P^{\mathfrak{M}'}$ .

mod:bas:red: prop:reduct **Proposition bas.2.** If an  $\mathcal{L}$ -**structure**  $\mathfrak{M}$  is a *reduct* of an  $\mathcal{L}'$ -**structure**  $\mathfrak{M}'$ , then for all  $\mathcal{L}$ -**sentences**  $\varphi$ ,

$$\mathfrak{M} \models \varphi \text{ iff } \mathfrak{M}' \models \varphi.$$

*Proof.* Exercise. □

**Problem bas.1.** Prove [Proposition bas.2](#).

**Definition bas.3.** When we have an  $\mathcal{L}$ -**structure**  $\mathfrak{M}$ , and  $\mathcal{L}' = \mathcal{L} \cup \{P\}$  is the expansion of  $\mathcal{L}$  obtained by adding a single  $n$ -place **predicate symbol**  $P$ , and  $R \subseteq |\mathfrak{M}|^n$  is an  $n$ -place relation, then we write  $(\mathfrak{M}, R)$  for the expansion  $\mathfrak{M}'$  of  $\mathfrak{M}$  with  $P^{\mathfrak{M}'} = R$ .

## Photo Credits

## Bibliography