

## syn.1 Unique Readability

lam:syn:unq:  
sec

We may wonder if for each term there is a unique way of forming it, and there is. For each lambda term there is only one way to construct and interpret it. In the following discussion, a *formation* is the procedure of constructing a term using the formation rules (one or several times) of ??.

lam:syn:unq:  
lem:term-start

**Lemma syn.1.** *A term starts with either a variable or a parenthesis.*

*Proof.* Something counts as a term only if it is constructed according to ??. If it is the result of ??, it must be a variable. If it is the result of ?? or ??, it starts with a parenthesis.  $\square$

lam:syn:unq:  
lem:app-start

**Lemma syn.2.** *The result of an application starts with either two parentheses or a parenthesis and a variable.*

*Proof.* If  $M$  is the result of an application, it is of the form  $(PQ)$ , so it begins with a parenthesis. Since  $P$  is a term, by **Lemma syn.1**, it begins either with a parenthesis or a variable.  $\square$

lam:syn:unq:  
lem:initial

**Lemma syn.3.** *No proper initial part of a term is itself a term.*

**Problem syn.1.** Prove **Lemma syn.3** by induction on the length of terms.

lam:syn:unq:  
prop:unq

**Proposition syn.4 (Unique Readability).** *There is a unique formation for each term. In other words, if a term  $M$  is formed by a formation, then it is the only formation that can form this term.*

*Proof.* We prove this by induction on the formation of terms.

1.  $M$  is of the form  $x$ , where  $x$  is some variable. Since the results of abstractions and applications always start with parentheses, they cannot have been used to construct  $M$ ; Thus, the formation of  $M$  must be a single step of ?????.
2.  $M$  is of the form  $(\lambda x. N)$ , where  $x$  is some variable and  $N$  is a term. It could not have been constructed according to ?????, because it is not a single variable. It is not the result of an application, by **Lemma syn.2**. Thus  $M$  can only be the result of an abstraction on  $N$ . By inductive hypothesis we know that formation of  $N$  is itself unique.
3.  $M$  is of the form  $(PQ)$ , where  $P$  and  $Q$  are terms. Since it starts with a parentheses, it cannot also be constructed by ?????. By **Lemma syn.1**,  $P$  cannot begin with  $\lambda$ , so  $(PQ)$  cannot be the result of an abstraction. Now suppose there were another way of constructing  $M$  by application, e.g., it is also of the form  $(P'Q')$ . Then  $P$  is a proper initial segment of  $P'$  (or vice versa), and this is impossible by **Lemma syn.3**. So  $P$  and  $Q$  are uniquely determined, and by inductive hypothesis we know that formations of  $P$  and  $Q$  is unique.  $\square$

A more readable paraphrase of the above proposition is as follows:

**Proposition syn.5.** *A term  $M$  can only be one of the following forms:*

1.  $x$ , where  $x$  is a variable uniquely determined by  $M$ .
2.  $(\lambda x. N)$ , where  $x$  is a variable and  $N$  is another term, both of which is uniquely determined by  $M$ .
3.  $(PQ)$ , where  $P$  and  $Q$  are two terms uniquely determined by  $M$ .

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## Bibliography