

## syn.1 Terms

lam:syn:trm: sec The terms of the lambda calculus are built up inductively from an infinite supply of variables  $v_0, v_1, \dots$ , the symbol “ $\lambda$ ”, and parentheses. We will use  $x, y, z, \dots$  to designate variables, and  $M, N, P, \dots$  to designate terms.

lam:syn:trm: defn:term **Definition syn.1 (Terms).** The set of *terms* of the lambda calculus is defined inductively by:

- lam:syn:trm: defn:term-var 1. If  $x$  is a variable, then  $x$  is a term.
- lam:syn:trm: defn:term-abs 2. If  $x$  is a variable and  $M$  is a term, then  $(\lambda x. M)$  is a term.
- lam:syn:trm: defn:term-app 3. If both  $M$  and  $N$  are terms, then  $(MN)$  is a term.

If a term  $(\lambda x. M)$  is formed according to (2) we say it is the result of an *abstraction*, and the  $x$  in  $\lambda x$  is called a *parameter*. A term  $(MN)$  formed according to (3) is the result of an *application*.

The terms defined above are fully parenthesized. This can get rather cumbersome, as the term  $(\lambda x. ((\lambda x. x)(\lambda x. (xx))))$  demonstrates. We will introduce conventions for avoiding parentheses. However, the official definition makes it easy to determine how a term is constructed according to **Definition syn.1**. For example, the last step of forming the term  $(\lambda x. ((\lambda x. x)(\lambda x. (xx))))$  must be abstraction where the *parameter* is  $x$ . It results by abstraction from the term  $((\lambda x. x)(\lambda x. (xx)))$ , which is an application of two terms. Each of these two terms is the result of an abstraction, and so on.

**Problem syn.1.** Describe the formation of  $(\lambda g. (\lambda x. (g(xx)))(\lambda x. (g(xx))))$ .

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## Bibliography