

int.1 Syntax of Intuitionistic Logic

int:int:syn:
sec

The syntax of intuitionistic logic is the same as that for propositional logic. In classical propositional logic it is possible to define connectives by others, e.g., one can define $\varphi \rightarrow \psi$ by $\neg\varphi \vee \psi$, or $\varphi \vee \psi$ by $\neg(\neg\varphi \wedge \neg\psi)$. Thus, presentations of classical logic often introduce some connectives as abbreviations for these definitions. This is not so in intuitionistic logic, with two exceptions: $\neg\varphi$ can be—and often is—defined as an abbreviation for $\varphi \rightarrow \perp$. Then, of course, \perp must not itself be defined! Also, $\varphi \leftrightarrow \psi$ can be defined, as in classical logic, as $(\varphi \rightarrow \psi) \wedge (\psi \rightarrow \varphi)$.

Formulas of propositional intuitionistic logic are built up from *propositional variables* and the propositional constant \perp using *logical connectives*. We have:

1. A denumerable set At_0 of *propositional variables* p_0, p_1, \dots
2. The propositional constant for *falsity* \perp .
3. The logical connectives: \wedge (conjunction), \vee (disjunction), \rightarrow (conditional)
4. Punctuation marks: $(,)$, and the comma.

int:int:syn:
defn:formulas

Definition int.1 (Formula). The set $\text{Frm}(\mathcal{L}_0)$ of *formulas* of propositional intuitionistic logic is defined inductively as follows:

1. \perp is an atomic *formula*.
2. Every *propositional variable* p_i is an atomic *formula*.
3. If φ and ψ are *formulas*, then $(\varphi \wedge \psi)$ is a *formula*.
4. If φ and ψ are *formulas*, then $(\varphi \vee \psi)$ is a *formula*.
5. If φ and ψ are *formulas*, then $(\varphi \rightarrow \psi)$ is a *formula*.
6. Nothing else is a *formula*.

In addition to the primitive connectives introduced above, we also use the following *defined* symbols: \neg (negation) and \leftrightarrow (*biconditional*). Formulas constructed using the defined operators are to be understood as follows:

1. $\neg\varphi$ abbreviates $\varphi \rightarrow \perp$.
2. $\varphi \leftrightarrow \psi$ abbreviates $(\varphi \rightarrow \psi) \wedge (\psi \rightarrow \varphi)$.

Although \neg is officially treated as an abbreviation, we will sometimes give explicit rules and clauses in definitions for \neg as if it were primitive. This is mostly so we can state practice problems.

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Bibliography