

seq.1 Derivations with Identity predicate

fol:seq:ide: Derivations with identity predicate require additional initial sequents and
sec inference rules.

Definition seq.1 (Initial sequents for =). If t is a closed term, then $\Rightarrow t = t$ is an initial sequent.

The rules for = are (t_1 and t_2 are closed terms):

$$\frac{t_1 = t_2, \Gamma \Rightarrow \Delta, \varphi(t_1)}{t_1 = t_2, \Gamma \Rightarrow \Delta, \varphi(t_2)} = \qquad \frac{t_1 = t_2, \Gamma \Rightarrow \Delta, \varphi(t_2)}{t_1 = t_2, \Gamma \Rightarrow \Delta, \varphi(t_1)} =$$

Example seq.2. If s and t are closed terms, then $s = t, \varphi(s) \vdash \varphi(t)$:

$$\frac{\frac{\varphi(s) \Rightarrow \varphi(s)}{s = t, \varphi(s) \Rightarrow \varphi(s)} \text{WL}}{s = t, \varphi(s) \Rightarrow \varphi(t)} =$$

This may be familiar as the principle of substitutability of identicals, or Leibniz' Law.

LK proves that = is symmetric and transitive:

$$\frac{\frac{\Rightarrow t_1 = t_1} {t_1 = t_2 \Rightarrow t_1 = t_1} \text{WL}}{t_1 = t_2 \Rightarrow t_2 = t_1} = \qquad \frac{\frac{t_1 = t_2 \Rightarrow t_1 = t_2} {t_2 = t_3, t_1 = t_2 \Rightarrow t_1 = t_2} \text{WL}}{t_2 = t_3, t_1 = t_2 \Rightarrow t_1 = t_3} = \frac{t_1 = t_2, t_2 = t_3 \Rightarrow t_1 = t_3} {t_1 = t_2, t_2 = t_3 \Rightarrow t_1 = t_3} \text{XL}$$

In the proof on the left, the formula $x = t_1$ is our $\varphi(x)$. On the right, we take $\varphi(x)$ to be $t_1 = x$.

Problem seq.1. Give derivations of the following sequents:

1. $\Rightarrow \forall x \forall y ((x = y \wedge \varphi(x)) \rightarrow \varphi(y))$
2. $\exists x \varphi(x) \wedge \forall y \forall z ((\varphi(y) \wedge \varphi(z)) \rightarrow y = z) \Rightarrow \exists x (\varphi(x) \wedge \forall y (\varphi(y) \rightarrow y = x))$

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Bibliography