

axd.1 Derivability and the Propositional Connectives

fol:axd:ppr:sec We establish that the **derivability** relation \vdash of axiomatic deduction is strong explanation enough to establish some basic facts involving the propositional connectives, such as that $\varphi \wedge \psi \vdash \varphi$ and $\varphi, \varphi \rightarrow \psi \vdash \psi$ (modus ponens). These facts are needed for the proof of the completeness theorem.

Proposition axd.1.

fol:axd:ppr:prop:provability-land
fol:axd:ppr:prop:provability-land-left
fol:axd:ppr:prop:provability-land-right

1. Both $\varphi \wedge \psi \vdash \varphi$ and $\varphi \wedge \psi \vdash \psi$
2. $\varphi, \psi \vdash \varphi \wedge \psi$.

Proof. 1. From ?? and ?? by modus ponens.

2. From ?? by two applications of modus ponens. □

Proposition axd.2.

fol:axd:ppr:prop:provability-lor

1. $\varphi \vee \psi, \neg\varphi, \neg\psi$ is inconsistent.
2. Both $\varphi \vdash \varphi \vee \psi$ and $\psi \vdash \varphi \vee \psi$.

Proof. 1. From ?? we get $\vdash \neg\varphi \rightarrow (\varphi \rightarrow \perp)$ and $\vdash \neg\psi \rightarrow (\psi \rightarrow \perp)$. So by the deduction theorem, we have $\{\neg\varphi\} \vdash \varphi \rightarrow \perp$ and $\{\neg\psi\} \vdash \psi \rightarrow \perp$. From ?? we get $\{\neg\varphi, \neg\psi\} \vdash (\varphi \vee \psi) \rightarrow \perp$. By the deduction theorem, $\{\varphi \vee \psi, \neg\varphi, \neg\psi\} \vdash \perp$.

2. From ?? and ?? by modus ponens. □

Proposition axd.3.

fol:axd:ppr:prop:provability-lif
fol:axd:ppr:prop:provability-lif-left
fol:axd:ppr:prop:provability-lif-right

1. $\varphi, \varphi \rightarrow \psi \vdash \psi$.
2. Both $\neg\varphi \vdash \varphi \rightarrow \psi$ and $\psi \vdash \varphi \rightarrow \psi$.

Proof. 1. We can **derive**:

1. φ HYP
2. $\varphi \rightarrow \psi$ HYP
3. ψ 1, 2, MP

2. By ?? and ?? and the deduction theorem, respectively. □

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Bibliography