## axd.1 Derivability and Consistency

fol:axd:prv: We will now establish a number of properties of the derivability relation. They are independently interesting, but each will play a role in the proof of the completeness theorem.

fol:axd:prv: **Proposition axd.1.** If  $\Gamma \vdash \varphi$  and  $\Gamma \cup \{\varphi\}$  is inconsistent, then  $\Gamma$  is inconprop:provability-contr sistent.

*Proof.* If  $\Gamma \cup \{\varphi\}$  is inconsistent, then  $\Gamma \cup \{\varphi\} \vdash \bot$ . By ??,  $\Gamma \vdash \psi$  for every  $\psi \in \Gamma$ . Since also  $\Gamma \vdash \varphi$  by hypothesis,  $\Gamma \vdash \psi$  for every  $\psi \in \Gamma \cup \{\varphi\}$ . By ??,  $\Gamma \vdash \bot$ , i.e.,  $\Gamma$  is inconsistent.

fol:axd:prv: **Proposition axd.2.**  $\Gamma \vdash \varphi$  iff  $\Gamma \cup \{\neg\varphi\}$  is inconsistent.

*Proof.* First suppose  $\Gamma \vdash \varphi$ . Then  $\Gamma \cup \{\neg \varphi\} \vdash \varphi$  by ??.  $\Gamma \cup \{\neg \varphi\} \vdash \neg \varphi$  by ??. We also have  $\vdash \neg \varphi \rightarrow (\varphi \rightarrow \bot)$  by ??. So by two applications of ??, we have  $\Gamma \cup \{\neg \varphi\} \vdash \bot$ .

Now assume  $\Gamma \cup \{\neg \varphi\}$  is inconsistent, i.e.,  $\Gamma \cup \{\neg \varphi\} \vdash \bot$ . By the deduction theorem,  $\Gamma \vdash \neg \varphi \rightarrow \bot$ .  $\Gamma \vdash (\neg \varphi \rightarrow \bot) \rightarrow \neg \neg \varphi$  by ??, so  $\Gamma \vdash \neg \neg \varphi$  by ??. Since  $\Gamma \vdash \neg \neg \varphi \rightarrow \varphi$  (??), we have  $\Gamma \vdash \varphi$  by ?? again.

**Problem axd.1.** Prove that  $\Gamma \vdash \neg \varphi$  iff  $\Gamma \cup \{\varphi\}$  is inconsistent.

fol:axd:prv: **Proposition axd.3.** If  $\Gamma \vdash \varphi$  and  $\neg \varphi \in \Gamma$ , then  $\Gamma$  is inconsistent.

*Proof.*  $\Gamma \vdash \neg \varphi \rightarrow (\varphi \rightarrow \bot)$  by **??**.  $\Gamma \vdash \bot$  by two applications of **??**.

fol:axd:prv: **Proposition axd.4.** If  $\Gamma \cup \{\varphi\}$  and  $\Gamma \cup \{\neg\varphi\}$  are both inconsistent, then  $\Gamma$  prop:provability-exhaustive is inconsistent.

Proof. Exercise.

Problem axd.2. Prove Proposition axd.4

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Bibliography