

## axd.1 Derivations with Identity predicate

fol:axd:ide:  
sec In order to accommodate = in **derivations**, we simply add new axiom schemas. The definition of **derivation** and  $\vdash$  remains the same, we just also allow the new axioms.

**Definition axd.1** (Axioms for **identity predicate**).

fol:axd:ide: 
$$t = t, \tag{1}$$

fol:axd:ide:  
ax:id1 
$$t_1 = t_2 \rightarrow (\psi(t_1) \rightarrow \psi(t_2)), \tag{2}$$

ax:id2

for any ground terms  $t, t_1, t_2$ .

fol:axd:ide:  
prop:sound **Proposition axd.2.** *The axioms eq. (1) and eq. (2) are valid.*

*Proof.* Exercise. □

**Problem axd.1.** Prove **Proposition axd.2**.

fol:axd:ide:  
prop:iden1 **Proposition axd.3.**  $\Gamma \vdash t = t$ , for any term  $t$  and set  $\Gamma$ .

fol:axd:ide:  
prop:iden2 **Proposition axd.4.** If  $\Gamma \vdash \varphi(t_1)$  and  $\Gamma \vdash t_1 = t_2$ , then  $\Gamma \vdash \varphi(t_2)$ .

*Proof.* The **formula**

$$(t_1 = t_2 \rightarrow (\varphi(t_1) \rightarrow \varphi(t_2)))$$

is an instance of **eq. (2)**. The conclusion follows by two applications of MP. □

## Photo Credits

## Bibliography