

## thy.1 The $s$ - $m$ - $n$ Theorem

cmp:thy:smn:  
sec The next theorem is known as the “ $s$ - $m$ - $n$  theorem,” for a reason that will explanation be clear in a moment. The hard part is understanding just what the theorem says; once you understand the statement, it will seem fairly obvious.

cmp:thy:smn:  
thm:s-m-n **Theorem thy.1.** *For each pair of natural numbers  $n$  and  $m$ , there is a primitive recursive function  $s_n^m$  such that for every sequence  $x, a_0, \dots, a_{m-1}, y_0, \dots, y_{n-1}$ , we have*

$$\varphi_{s_n^m(x, a_0, \dots, a_{m-1})}^n(y_0, \dots, y_{n-1}) \simeq \varphi_x^{m+n}(a_0, \dots, a_{m-1}, y_0, \dots, y_{n-1}).$$

It is helpful to think of  $s_n^m$  as acting on *programs*. That is,  $s_n^m$  takes a explanation program,  $x$ , for an  $(m+n)$ -ary function, as well as fixed inputs  $a_0, \dots, a_{m-1}$ ; and it returns a program,  $s_n^m(x, a_0, \dots, a_{m-1})$ , for the  $n$ -ary function of the remaining arguments. If you think of  $x$  as the description of a Turing machine, then  $s_n^m(x, a_0, \dots, a_{m-1})$  is the Turing machine that, on input  $y_0, \dots, y_{n-1}$ , prepends  $a_0, \dots, a_{m-1}$  to the input string, and runs  $x$ . Each  $s_n^m$  is then just a primitive recursive function that finds a code for the appropriate Turing machine.

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## Bibliography